# Efficient Analysis of Population Protocols and Chemical Reaction Networks 

Doctoral Defense

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## Introduction: Motivation



## Introduction: Models

"Population protocols and chemical reaction networks are formal models in which many simple entities interact resulting in a hard to analyse system."

## Introduction: Use Cases



Self-Organizing Systems

## Introduction: Thesis Overview

## Goal of Thesis:

Enable efficient (or "practical") analysis of these models via:

1. Theoretical results
2. New efficient analysis methods
3. Easily accessible tools

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## Results of Thesis:

- Efficient verification of population protocols
$\rightarrow$ in this talk: focus on tool Peregrine
- Synthesis of efficient population protocols
$\rightarrow$ not in this talk
- Efficient transient analysis of chemical reaction networks
$\rightarrow$ in this talk: focus on segmental simulation idea

Population Protocols


- Anonymous mobile agents with very few resources

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At least as many blue birds as red birds?


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Protocol:

- 4 states: blue/red, large/small
- Two large birds of different colors become small
- Large birds convert small birds to their color
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Correctness properties:

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$$
\begin{aligned}
& (\dot{\phi} \geq \dot{\phi}) \Longrightarrow \text { PG }(\dot{\phi}+\boldsymbol{\phi}=0) \\
& (\dot{\phi})<\dot{\phi}) \Longrightarrow \text { PG }(\dot{\phi}+\boldsymbol{\phi}=0)
\end{aligned}
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"Birds converge to color of majority."

## Population Protocols: Verification Idea

## Definition: Configuration

A vector describing the number of agents per state:

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Correctness proofs are typically structured in stages that trap the system in progressively more constrained subsets of configurations.

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Correctness proofs are typically structured in stages that trap the system in progressively more constrained subsets of configurations.
$\rightarrow$ Idea: Formalize structure as stage graph!

## Population Protocols: Stage Graphs

## Definition: Stage Graph

A stage graph for $\varphi_{\text {pere }} \Rightarrow F G \varphi_{\text {post }}$ is a finite DAG satisfying:

1. Each node is an inductive set of configurations, called stage. "It is impossible to leave stages."

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3. In stages with successors, executions enters a successor with prob. 1. "Stages lead to their successors."
4. In stages without successors, all configurations satisfy $\varphi_{\text {post }}$. "Finally, the postcondition holds forever."

Example:


## Population Protocols: Results

## Theory

A population protocol is correct if and only if there is a stage graph proving it.

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## Practice

- Most systems have small stage graphs
- Most systems make progress by "killing" transitions
$\rightarrow$ Use heuristics to efficiently construct Presburger stage graphs


## Population Protocols: Utility of Stage Graphs

Stage graphs:

- Can be efficiently constructed for most protocols
- Certify liveness properties
- Are independently checkable
- Explain how the protocol works
- Give speed guarantees
- Help to find counter examples


## 0120 <br> 

Try it in Peregrine!

## Chemical Reaction Networks

## Chemical Reaction Networks

## Example: Viral Infection

| Species: | RNA, DNA, V, P |  |
| :--- | :--- | :--- |
| Initial state: | $(1 \times \mathrm{RNA})$ |  |
| End time: | 200 s |  |
| Reactions: | DNA $+\mathrm{P} \xrightarrow{0.00001} \mathrm{~V}$ | $\mathrm{RNA} \xrightarrow{0.25} \emptyset$ |
|  | RNA $\xrightarrow{1000}$ RNA +P | $\mathrm{P} \xrightarrow{2} \emptyset$ |
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Similar to population protocols but:

- Variable number of molecules
- Continuous time (,i.e., behaves like CTMC not DTMC)
- Focus on modeling systems (instead of designing them)


## CRN: Transient Analysis

How does the system evolve?

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Transient analysis:

- Hard because of complex dynamics, state-space explosion, stochasticity, stiffness, multimodality, ...
- Two approaches:
- Direct (numerical)
- Indirect (using many trajectories)


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## This work

- Goal: Efficiently compute many simulations
- Idea: Use memoization!


## CRN: Simulation

Gillespie's stochastic simulation algorithm (SSA) [9]

- Sample one reaction at a time


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$$
t=0 s
$$

Sinit

Start in initial state.

## CRN: Simulation

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- Sample one reaction at a time


Compute rate of all reactions.

## CRN: Simulation

## Gillespie's stochastic simulation algorithm (SSA) [9]

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Time until the next reaction: $\Delta t \sim \operatorname{EXP}(0.5+2+1.5)$
Probability of reactions: $\frac{0.5}{4}, \frac{2}{4}, \frac{1.5}{4}$

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Gillespie's stochastic simulation algorithm (SSA) [9]

- Sample one reaction at a time

$$
t=0.18 \mathrm{~s}
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Apply change to both time and state, then repeat.

## CRN: Simulation

## Gillespie's stochastic simulation algorithm (SSA) [9]

- Sample one reaction at a time


Time until the next reaction: $\Delta t \sim \operatorname{EXP}(1+2+2)$
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## Gillespie's stochastic simulation algorithm (SSA) [9]

- Sample one reaction at a time


Time until the next reaction: $\Delta t \sim E X P(4+2+3+1)$
Probability of reactions: $\frac{4}{10}, \frac{2}{10}, \frac{3}{10}, \frac{1}{10}$

## CRN: Simulation

## Gillespie's stochastic simulation algorithm (SSA) [9]

- Sample one reaction at a time
- May take a long time



## CRN: Segmental Simulation

Precompute $k$ short trajectories (called segments) for each state.
$\rightarrow$ Simulate by sampling segments instead of single reactions.


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- much faster!
- Problem: many states $\rightarrow$ too inefficient


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## CRN: Abstraction-Based Segmental Simulation

Only precompute $k$ segments for each representative.


MEMORY

$$
(k=3)
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Segments end when they leave the abstract state.
$\rightarrow$ Intuition: "significant change"

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To save memory: Work with summaries instead of segments.

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Lazy: Do not precompute but fill memory on-the-fly!

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## CRN: Conclusion

Segmental simulation:

- Reuses previous simulations to generate new ones
- Is an approximate simulation technique
- Accurately captures dynamics of most systems
- Speeds up transient analysis (up to 4000x faster)
- Efficiently predict the behavior of CRNs


Try it in SAQuaiA!

Thank you


## My Publications

## First Author:

- Checking Qualitative Liveness Properties of Replicated Systems with Stochastic Scheduling (CAV'20) [3]
- Peregrine 2.0: Explaining Correctness of Population Protocols Through Stage Graphs (ATVA'20) [8]
- Abstraction-Based Segmental Simulation of Chemical Reaction Networks (CMSB'22) [10]


## Non-first Author:

- Succinct Population Protocols for Presburger Arithmetic (STACS'20) [2]
- Fast and Succinct Population Protocols for Presburger Arithmetic (SAND'22) [5]


## My Publications if

## Others (not part of thesis):

- Automata Tutor v3 (CAV20) [7]
- Decision Power of Weak Asynchronous Models of Distributed Computing (PODC'21) [4]
- Fast and succinct population protocols for Presburger arithmetic (Journal of Comp. and Sys. Sciences 2023) [6]


## Population Protocol: Synthesis

## Our result

For every quantifier-free Presburger formula $\varphi$ there is an population protocol which

- has $\mathcal{O}(\operatorname{POLY}(|\varphi|))$ states,
$\rightarrow$ succinct
- for $n$ agents stabilizes in $\mathcal{O}\left(n^{2} \log n\right)$ expected interactions, and $\rightarrow$ fast
- can be constructed efficiently.



## Population Protocol: Synthesis Idea

## Idea:

- Binary value representation (instead of unary)
- Agents participate in computation of only one atomic formula (instead of in all)


## Population Computer

Like population protocol but easier to design because:

- k-way transitions
- Helper agents
- More flexible output definition

Population Protocol: Synthesis (Fast and Succinct)


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